

Formal Representation and Automated Transformation of Geometric Statements

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Outline

- 1 Motivation
- 2 Geometry Programming Language
- 3 Geometric Statement Simplification
- 4 Implementation
- 5 Conclusion and Future Work

Start with Geometry Software

Geometry problems (drawing or proving) are specified by applying similar (or same) concepts which are implemented differently in these systems.

Table: Constructive style

Cinderella	GeoGebra
Perpendicular(a;A)	PerpendicularLine[A,a]
Circumcircle(A;B;C)	Circle[A,B,C]
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parallel(A,B,C,D)	parallel A B C D

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The constructions and predicates can be viewed as **concepts**.

Standardizing Problem Specifications

It is needed to **standardize the formats of specifications** so that the same specified problems can be processed by different geometry software systems via specific interfaces.

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- **Intergeo** project offers a common file format for specifying dynamic diagrams. However, the format only works for constructive style.
- **GeoCode** is a generic proof scheme standard providing routine codes that can be interfaced with different CAS or provers for proving and DGS for drawing.

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It is needed to **standardize macro constructions** so that one can specify problems in terms of customized concepts by defining macros.

Related work

- **GEOTHER** provides a standard form for specifying the entries contained in the predicates routines. However, defined predicates are independent with each other.
- **GeoCode** provides the facility for users to define new functions in terms of existed functions. However, these functions are defined only in the constructive style.

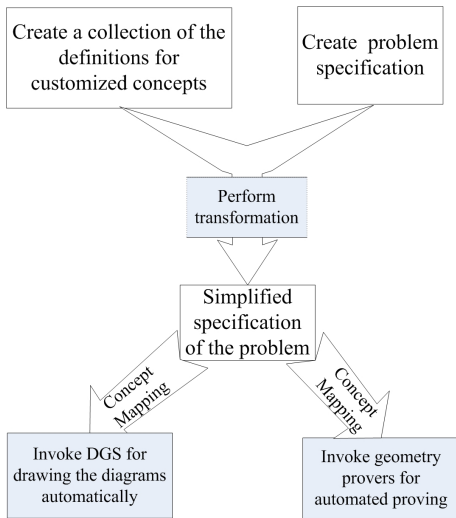
Objectives

- A general **geometry programming language** is needed in which one can **easily** and **naturally** define geometric concepts and specify problems in terms of the customized concepts (for both constructive and constraint type).

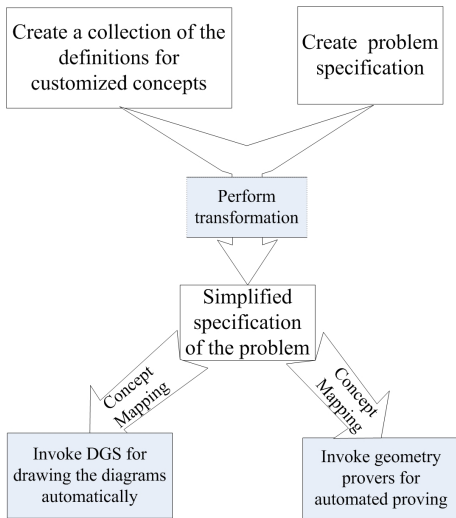
Objectives

- A general **geometry programming language** is needed in which one can **easily** and **naturally** define geometric concepts and specify problems in terms of the customized concepts (for both constructive and constraint type).
- The facility is needed for **transforming** the specified problems into the ones that target systems can identify and manipulate via specific interfaces.

Idea



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[▶ see Demo](#)

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Concept Symbols

- **Customized concepts:** point, line, intersection, midpoint, area, etc.

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- **Customized concepts:** point, line, intersection, midpoint, area, etc.
- **Built-in concepts:**
 - **Constants:** 0 , π , etc.
 - **Pointers (labels):** A , B , l , etc.
 - **Types:** Point, Line, Segment, Length, Degree, Number, Boolean, etc.
 - **Algebra concepts:** *times*, *plus*, *sin*, *squre*, etc.
 - **Set concepts:** *list*, *choose*, *ismember*, etc.
 - **Logic concepts:** *and*, *or*, *not*.

Formalization of Geometric Concepts

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- **Entity concepts:**
 - **Geometric objects:** $\text{intersection}(l :: \text{Line}, m :: \text{Line})$,
 $\text{perpendicularline}(A :: \text{Point}, l :: \text{Line})$,
 $\text{circumcenter}(\text{triangle}(A :: \text{Point}, B :: \text{Point}, C :: \text{Point}))$, etc.
 - **Geometric quantities:** $\text{length}(\text{segment}(A :: \text{Point}, B :: \text{Point}))$,
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- **Boolean concepts:**
 - **Geometric relations:** $\text{parallel}(l::\text{Line},m::\text{Line})$, $\text{isin}(A::\text{Point},o::\text{Circle})$,
 $\text{tangent}(o::\text{Circle},p::\text{Circle})$ etc.
 - **Quantity relations:** $\text{lt}(a::\text{Length},b::\text{Length})$,
 $\text{equal}(c::\text{Degree},d::\text{Degree})$, etc.

Constructing Geometric Clauses

Clauses are constructed by using **instances** (of concepts).

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- **Reference clauses:** $A := \text{point}()$, $P := \text{intersection}(l, m)$, etc.
- **Boolean clauses:** $\text{perpendicular}(l, m)$, $\text{incident}(A, l)$, etc.
- **Compound clauses:**
 - **Nesting:** $\text{collinear}(\text{foot}(D, \text{line}(A, B)), \text{foot}(D, \text{line}(A, C)), \text{foot}(D, \text{line}(B, C)))$;
 - **Give:** $\text{give}(\text{triangle}(A, B, C))$;
 - **Configuration:** $\text{configuration}(E := \text{intersection}(\text{line}(A, B), \text{line}(C, D)), F := \text{intersection}(\text{line}(A, C), \text{line}(B, D)))$;
 - **Declare:** $\text{declare}(A :: \text{Point}, B :: \text{Point}, l :: \text{Line})$;
 - **Logic:** $\text{and}(\text{parallel}(l, m), \text{incident}(A, l))$;
 - **List:** $\{A; B; C\}, \{\text{point}(); \text{point}(); \text{midpoint}(\text{segment}(A, B))\}$;
 - **Set:** $\text{choosediff}(A; B; C, 2)$;
 - **Algebra:** $\text{times}(2, \text{length}(\text{segment}(A, B)))$.

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- Definition(completequadrilateral($A::\text{Point}, B::\text{Point}, C::\text{Point}, D::\text{Point}, E::\text{Point}, F::\text{Point}$), [configuration($E:=\text{intersection}(\text{line}(A, B), \text{line}(C, D))$), $F:=\text{intersection}(\text{line}(A, C), \text{line}(B, D))$)], null)

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- Definition(diagonal(completequadrilateral($A::\text{Point}, B::\text{Point}, C::\text{Point}, D::\text{Point}, E::\text{Point}, F::\text{Point}$)), {[segment(A, D)]; [segment(B, C)]; [segment(E, F)]}, null)

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For example,

- `Problem(Simson, Theorem, assume(A:=point(), B:=point(), C:=point(), D:=point(), incident(D, circumcircle(triangle(A, B, C))))), show(collinear(foot(D, line(A, B)), foot(D, line(A, C)), foot(D, line(B, C))))`

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- Problem(Simson, Theorem, assume($A := \text{point}()$, $B := \text{point}()$, $C := \text{point}()$, $D := \text{point}()$, incident(D , circumcircle(triangle(A, B, C)))), show(collinear(foot(D , line(A, B)), foot(D , line(A, C)), foot(D , line(B, C))))))
- Problem(Pappus, Theorem, assume(declare($C :: \text{Point}$, $F :: \text{Point}$, $P :: \text{Point}$, $Q :: \text{Point}$, $R :: \text{Point}$), $A := \text{point}()$, $B := \text{point}()$, $D := \text{point}()$, $E := \text{point}()$), give(Pappus($A, B, C, D, E, F, P, Q, R$)), show(collinear(P, Q, R)))

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$$\triangleright \text{foot}(D, \text{line}(E, F)) \xrightarrow[\text{substitution}]{Def_1} \text{foot}(D, \text{line}(E, F)) \xrightarrow[\text{substitution}]{Def_2}$$

$$[\text{intersection}(\text{perpendicularline}(D, \text{line}(E, F)), \text{line}(E, F))]$$

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We adopt **eager (inner most)** strategy to deal with the nesting cases.

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show(collinear(foot(D,line(A,B)),foot(D,line(A,C)),foot(D,line(B,C))))))
```

$$\xrightarrow[\text{simplification}]{\text{definitions}}$$

```
Problem(Simson,Theorem,assume(declare(var_0::Point,var_1::Point,var_2::Line,
var_3::Point,var_4::Line,var_5::Point,var_6::Line,var_7::Point),
A:=point(),B:=point(),C:=point(),D:=point(),equal(distance(var_0,D),distance
(var_0,var_1)),equal(distance(var_0,var_1),distance(var_0,A)),equal(distance
(var_0,A),distance(var_0,B)),equal(distance(var_0,A),distance(var_0,C ),...),
show(incident(var_3,line(var_5,var_7))))))
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Type for instance

```
Type(foot(D,line(A,B))) = foot(point(),line(point()),point())
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```
Type(intersection(l,line(C,D))) =
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intersection(perpendicularline(point()),line(point()),point()),line(point()),point())
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Type for concept

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Type(foot(A::Point,l::Line)) = foot(Point,Line)
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Generally, type for instance is not equal to type for concept. How to match them?

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Geometry definitions indicate the order of types. We define **type upgrade** to match types.

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Example

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$\text{intersection}(\text{perpendicularline}(\text{point}(), \text{line}(\text{point}(), \text{point}())), \text{line}(\text{point}(), \text{point}()))$
 $< \text{intersection}(\text{perpendicularline}(\text{Point}, \text{Line}), \text{Line}) < \text{intersection}(\text{Line}, \text{Line})$

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`intersection(perpendicularline(point(),line(point(),point())),line(point(),point()))`
 < `intersection(perpendicularline(Point,Line),Line)` < `intersection(Line,Line)`

Type Matching Rule

Let I and C be an instance and a concept, if $\text{Type}(I) \leq \text{Type}(C)$, then the definition of C can be used to simplify instance I .

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[*I* **where** constraint **context** configuration **with** nondegeneracyCondition]

The simplified instances will be normalized into this form at each step of simplification process.

Analysis

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The process terminates only if there is no loop in the type structure determined by the definitions.

Usability

The simplified problem specifications can be interfaced with Geometry software systems.

More Demo

- **Reuse** definitions and problem specifications.

▶ Pappus,completeQuadrilateral,197,198

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- Dealing with multiple returns.

▶ example90, 180

Outline

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- 3 Geometric Statement Simplification
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- 5 Conclusion and Future Work

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We have presented a geometry programming language for specifying geometric concepts, definitions, and problems.

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- encoded **easily** and **naturally**;

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The specifications are

- encoded **easily** and **naturally**;
- used in both **constraint** and **constructive** cases;
- transformed into ones that can be **interfaced** with available geometry software systems.

Future Work

The geometry programming language is still at a preliminary stage. The following problems should be considered further.

- prove the correctness of transformation;
- transform the specifications in this language into natural language and the other way round;
- transform the specifications in this language into algebraic counterparts and interface with CAS.

Geo* - Geometry on Computer

Overview Research Team Publications Consortium Related Links Contact

Geo* - Geometry on Computer

The Geo* project attempts to bring the contents of traditional geometry to electronic form and to make geometric computation, reasoning, drawing, and knowledge management dynamic, automatic, or interactive on computer.

Current research in this project focuses on the

- identification, formalization, representation, and creation of geometric knowledge data and objects;
- design, implementation, and analysis of algorithms and software tools for geometric computation, reasoning, data processing, and diagram generation;
- development of methodologies and systems for geometric knowledge presentation and management;
- design and implementation of geometric specification and programming languages.

Project

- Data
- e-Text
- Draw
- Prove
- Lounge
- 中文

Welcome to visit our project home at <http://geo.cc4cm.org/>

Thanks!